Suppose you want to develop a data structure with:

- Insert(x): insert element x.
- Search(k): search for element with key k.
- **Delete**(x): delete element referenced by pointer x.
- ▶ find-by-rank(ℓ): return the ℓ -th element; return "error" if the data-structure contains less than ℓ elements.

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- verify/show how the additional information can be maintained for the basic modifying operations on the underlying structure.
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Goal: Design a data-structure that supports insert, delete, search, and find-by-rank in time $O(\log n)$.

- 1. We choose a red-black tree as the underlying data-structure.
- **2.** We store in each node v the size of the sub-tree rooted at v.
- 3. We need to be able to update the size-field in each node without asymptotically affecting the running time of insert, delete, and search. We come back to this step later...



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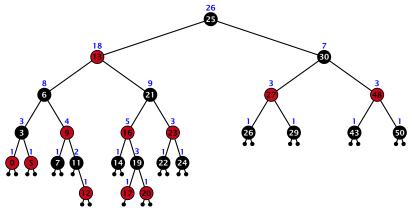
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4. How does find-by-rank work?
Find-by-rank(k) = Select(root,k) with

```
Algorithm 7 Select(x, i)
```

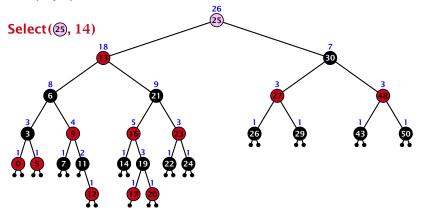
- 1: **if** x = null **then return** error
- 2: **if** left[x] \neq null **then** $r \leftarrow$ left[x]. size +1 **else** $r \leftarrow 1$
- 3: if i = r then return x
- 4: if i < r then
- 5: **return** Select(left[x], i)
- 6: else
- 7: **return** Select(right[x], i r)





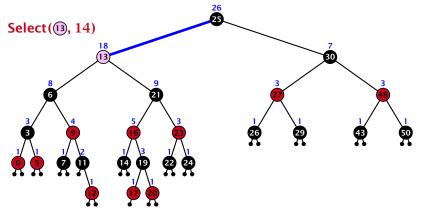
- decide whether you have to proceed into the left or right sub-tree
- adjust the rank that you are searching for if you go right





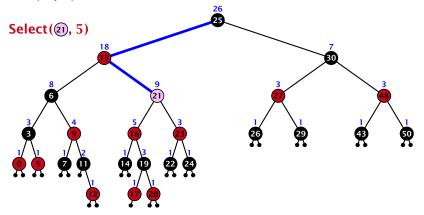
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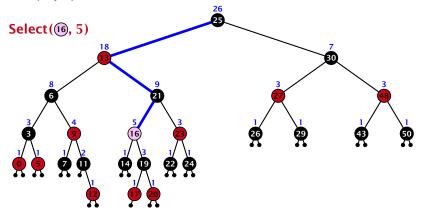
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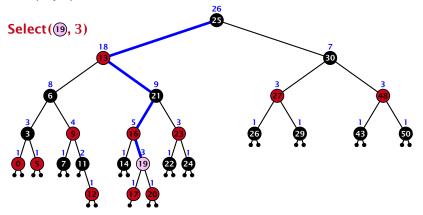
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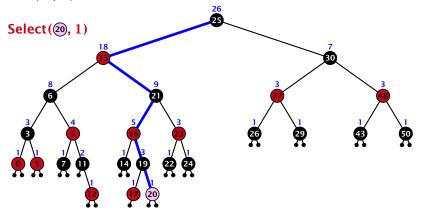
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3. How do we maintain information?

Search(k): Nothing to do.

Insert(x): When going down the search path increase the size field for each visited node. Maintain the size field during rotations.



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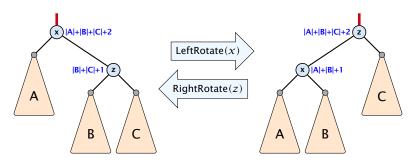
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Rotations

The only operation during the fix-up procedure that alters the tree and requires an update of the size-field:



The nodes x and z are the only nodes changing their size-fields.

The new size-fields can be computed locally from the size-fields of the children.



